

## A Light-Weighted Data Trust Model in WSN

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### Abstract

*WSN is often deployed in unattended or even hostile environments. Therefore, providing security in WSN is a major requirement for acceptance and deployment of WSN. Furthermore, establishing trust in a clustered environment can provide numerous advantages. We proposed a light-weighted trust model which consider data aggregation and detailed the data in parameters to depend against attacks. With comparing our model with LDTS and Model using Trust Matrix, we conclude that our model has implemented a trade-off between detection rate and communication consumption.*

**Keywords:** *direct trust, similarity, trust matrix, WSN, indirect trust*

### 1. Introduction

Applications ranging from health, home, environmental to military and defense make use of sensor nodes for collection of appropriate data. The sensor nodes comprising of data collecting, processing, and transmitting units are very small in size and can be densely deployed owing to their low cost [4]. Cluster WSN such as LEACH is broadly used. Clustering algorithms can effectively improve network scalability and throughput. Using clustering algorithms, nodes are grouped into clusters, and within each cluster, a node with strong computing power is elected as a cluster head (CH). CHs together form a higher-level backbone network. After several recursive iterations, a clustering algorithm constructs a multi-level WSN structure [5].

However, WSN is often deployed in unattended or even hostile environments. The wireless and resource-constraint nature of a sensor network makes it an ideal medium for attackers to do any kinds of vicious things. Therefore, providing security in WSN is a major requirement for acceptance and deployment of WSN [6]. Establishing trust in a clustered environment provides numerous advantages, such as enabling a CH to detect faulty or malicious nodes within a cluster. In the case of multi-hop clustering, a trust system aids in obtain correct data aggregation.

The rest of the paper is organized as follows. The models and definitions are proposed in Section 3. The detailed trust model is depicted in Section 4. The comparison and evaluation of our trust model with other models are given in Sections 5. The related work and our conclusions are presented in Sections 2 and 6.

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## 2. Related Work

Research on trust management systems for WSN received considerable attention from scholars. A number of studies have proposed such systems for WSNs. However, these systems suffer from various limitations such as the incapability to meet the resource constraint requirements of the WSNs, more specifically, for the large-scale WSN. Recently, very few trust management systems have been proposed for clustered WSNs, such as GTMS [1], RFSN [2], HATWA [7] and [3]. To our best knowledge, a universal trust system designed for clustered WSNs to achieve light-weight remains lacking.

In Group based Trust Management Scheme [1], the authors proposed a new light weight trust management scheme for WSN. It works with two different topologies: intragroup and intergroup, where distributed trust management and centralized trust management is adopted respectively. And the trust states are represented as Trusted, Untrusted and Uncertain respectively. The advantage of the scheme is that, it evaluates the trust for the group of nodes rather than a single node in the cluster. However, GTMS relies on a broadcast-based strategy to collect feedback from the CMs of a cluster, which requires a significant amount of resources and power.

In Reputation based Framework for Sensor Networks [2], each node in the network maintains the reputation for neighboring nodes. The values of the trust are evaluated on the basis of that reputation and Bayesian formulation was used for representing the values of reputation. According to this scheme, each node in the sensor network develops a reputation for other node in the network by making direct observations. The disadvantage of this model is that, the mobility makes there is no enough interaction with the nodes and hence the reputation information will not be stabilized.

HATWA proposes a Heuristic Approach based Trust Worthy Architecture for WSN that considers the challenges of the system and focus on the collaborative mechanism for trust evaluation and maintenance. The proposed Architecture could also be capable of fulfilling critical security, reliability, mobility and performance requirements for reliable communication while being readily adaptable to different applications. But it introduced an additional network monitoring node to record and calculate the whole trust during four stages. And data aggregation is not included in the four stages. They also assume the monitor node is no fault.

In a Fault-Event Detection Model Using Trust Matrix in WSN (DMUTM) [3], the author proposed a method of fault and event detection using trust model in WSN based on similarity matrix. They used similarity matrix which is based on data aggregation distinguish groups from each other in one cluster to detect fault. The trust was calculated by cluster head either directly or indirectly. When in indirectly case, the head calculated the trust by transitivity algorithm. However, the trust transitivity required a high complexity which leads to amount of power consumption.

In [8], Xiao proposed a trust system LDTS for WSNs, which employ clustering algorithms. First, a lightweight trust decision-making scheme is proposed based on the nodes' identities in the clustered WSNs. Then a dependability-enhanced trust evaluating approach is defined for co-operations between CHs. Moreover, a self-adaptive weighted method is defined for trust aggregation at CH level. But the method focuses on transmit process but not considers data property in the network. So it can only depend against Garnished attack and bad mouthing attack.

So it is necessary to build a light-weighted trust model which consider data aggregation and detailed the data in parameters to depend against more attacks. The contribution of our work is:

- 1) Create a light-weighted trust model based on our former work.
- 2) Parameterize the sensing data according to three factors.
- 3) Compare our model with LDTS and Model using Trust Matrix.

### 3. Models and Definitions

#### 3.1. Network Model

WSN in a two dimensional plane with  $n$  sensors, denoted by a set  $N = (n_1, n_2, \dots, n_n)$ , where  $n_i$  is the  $i$ th sensor. These sensors are placed in an area and the transmission radius is  $r_s$ . Each node maintains its ID, sensing data and location. In such a network, we use LEACH protocol to create clusters illustrated in Figure 1. A node in the clustered WSN model can be identified as a CH, or a CM (See Figure. 1). Members of a cluster can communicate with their CH directly. A CH can forward the aggregated data to the central BS through other CHs.

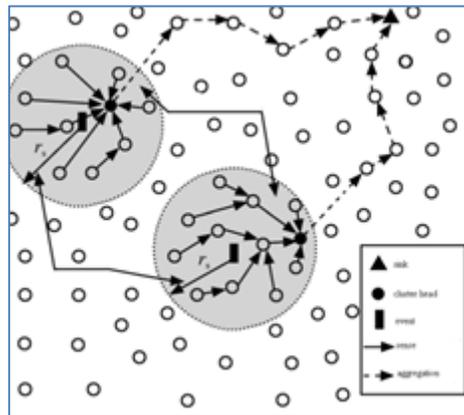


Figure 1. Model of Network

#### 3.2. Trust Model

Trust models are classified into two categories that are node trust models and data trust models [6].

A data trust model is proposed to distinguish forged data of illegal nodes from innocent data of legal nodes. Sensor nodes evaluate trustworthiness of their neighbor nodes by cross checking the neighbor nodes' redundant sensing data with their own result. The trust value is calculated through a light-weighted method, and the data considering is a structure composed of three parameters: the consistency value of sensing data, the communication ability and the remained lifetime of a node. After the trust assertion, inconsistent data from malicious or compromised nodes can be detected.

#### 3.3. Definition of Trust Matrix

When consider a cluster, we get a  $G = (V, E, s)$  consists of vertexes  $V$ , edges  $E$  and similarity weight  $s$ . Each vertex is a node and each edge is the connection of two neighbors. We compute the similarity among sensor nodes as formula (1) where node  $i$  and node  $j$  is adjacent in location.  $X$  is the sensing data of node. If  $s_{i,j} > 0.9$ , we set new  $s_{i,j}$  as 1, otherwise as 0.

$$s_{i,j} = \frac{X_i X_j}{X_i^2 + X_j^2 - X_i X_j} \quad (1)$$

We consider a window of time  $\Delta t$ . Thus, as time elapses, the window deletes old experiences but adds newer experiences. The trust value between two nodes can be calculated according to formula (2).

$$T_{x,y}(\Delta t) = \left\lceil \left( \frac{10 * s_{x,y}(\Delta t)}{s_{x,y}(\Delta t) + d_{x,y}(\Delta t)} \right) \left( \frac{1}{\sqrt{d_{x,y}(\Delta t)}} \right) \right\rceil \quad (2)$$

$\lceil \cdot \rceil$  is the nearest integer function.  $s_{x,y}(\Delta t)$  is the total number of similar data comparison of node x with y in  $\Delta t$  time, and  $d_{x,y}(\Delta t)$  is the total number of dissimilar data comparison. Specially, if  $d_{x,y}(\Delta t) = 0$ , we set  $T_{x,y}(\Delta t) = 10$ .

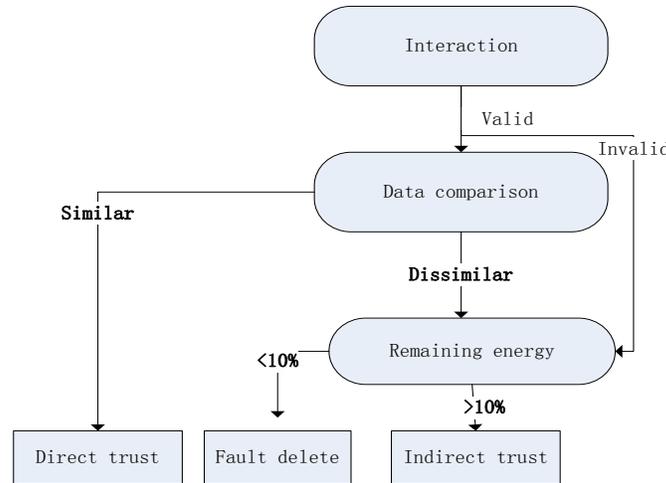
The cluster head will periodically broadcast the request packet within the cluster. In response, all CMs in the cluster will forward their data values to CH. Then, CH will maintain these values in a matrix as shown below where the real number is the similarity of node i for node j and 1 is a default value presenting the similarity of the node for itself.

$$\begin{matrix} s_{1,1} & \dots & s_{1,n} \\ \vdots & \ddots & \vdots \\ s_{n,1} & \dots & s_{n,n} \end{matrix} \quad (3)$$

## 4. A Light-weighted Trust Model

### 4.1. Calculate Direct Trust

A CM's trust value can be calculated by direct and indirect observation. Direct trust is evaluated by the number of successful and unsuccessful interactions, similar or dissimilar data comparison. In this work, interaction refers to the cooperation of two CMs and comparison refers to data aggregation. Indirect trust is evaluated by aid of similarity matrix in CH. That is, if node x wants to calculate the trust value for node y, first it checks whether it has a valid interaction with y during a specific time interval. If a past valid interaction record exists, then it compares its data value with y. Otherwise, if remain energy is less than ten percent, it will send a request to its CH. The model considers the consistency value of sensing data, the communication ability and the remained lifetime of a node. The process can be depicted in Figure 2.



**Figure 2. Process of the Model**

In the first stage of interaction, we use LDTS model in [8]. If the interaction  $I_{x,y}$  is more than 5, it is regarded as valid. Then start second stage to calculate data comparison using formula (2). If data similarity is more than 5, the direct trust is as formula (4). Otherwise, check the remaining energy to decide whether to calculated the indirect trust or assert the node is fault to delete from the network.

$$IT_{x,y} = \frac{I_{x,y} * S_{x,y}}{10} \quad (4)$$

If the interaction  $I_{x,y}$  is less than 5, it is regarded as invalid. Then we check the remaining energy to do the same work as above.

#### 4.2. Calculate Indirect Trust

When enter the stage of calculating indirect trust, the node request to CH for a feedback that if there are some nodes that are not its neighbors have the similar data as it sensing. We use the beta probability density functions to compute the indirect trust as formula (5) based on matrix (4).

$$DT_{ch,y} = 10 * \frac{s_{x,y} + 1}{s_{x,y} + d_{x,y} + 2} \quad (4)$$

Here,  $s_{x,y}$  denotes the number of similar data to node y except itself and  $d_{x,y}$  denotes the number of dissimilar data to node y at one moment. For example, as shown in Figure 4 which is deduced from Figure 3, we want to calculate indirect trust of node 1. The value is a real number of 6.7.

1	0.9	0	0.95	0.98
0.9	1	0.1	0	0.92
0	0.1	1	0	0.1
0.95	0	0	1	0.9
0.98	0.92	0.1	0.9	1

**Figure 3. Similarity Matrix**

$$\begin{bmatrix} 1 & 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 1 & 1 \\ 1 & 1 & 0 & 1 & 1 \end{bmatrix}$$

**Figure 4. Trust Matrix**

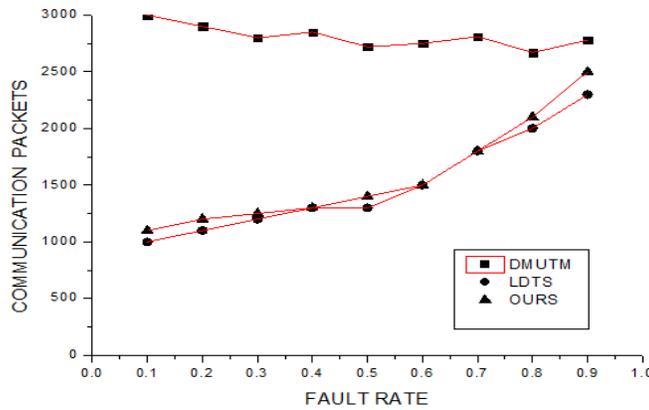
**5. Evaluation**

Our experiment uses ns3 to design. Fifty sensor nodes are distributed in a space of  $500 \times 700$ , and the communication radius is set as 60. Each node has two to five neighbors in the experiment and the node's location is already known. The detailed value is shown in Table 1.

**Table 1. Values in Evaluation**

Symbol	Description	values
N	Number of nodes	50
n	Number of CMs in a cluster	6-8
m	Number of CM's neighbors	4-6

We only consider the communication overhead with ignoring calculation cost. It is also assumed that the route is reliable without considering the case of route failure. We compare the communication consumption and error detection rate for simulation to LDTS and DMUTM.



**Figure 5. Comparison of Communication Consumption**

It is shown in Figure 5 that, when data error rate is changed, DMUTM maintains an average of packets by 2800. But LDTS and ours algorithm gets an increased average of packets as the data error rate growing. This is due to the calculating of indirect trust which will consume more communication. While for LDTS, it only consider interaction, so the probability of calculating indirect trust is less than ours since our method considers either the interaction and data similarity.

Except for energy consumption, error detection rate is another important merit to measure a trust algorithm. We define error detection rate as  $f_s/f$ , where  $f_s$  is the number of fault nodes that have been detected and  $f$  is the total number of fault nodes.

Simulation result shown in Figure 6 indicates that the detection rate of DMUTM is higher than the other two methods because it handles all cases with indirect trust. And the higher

detection rate is an exchange for communication consumption. LDTS has a lower detection rate than ours since it omit the data fault. Our model implements a balance between detection rate and communication consumption.

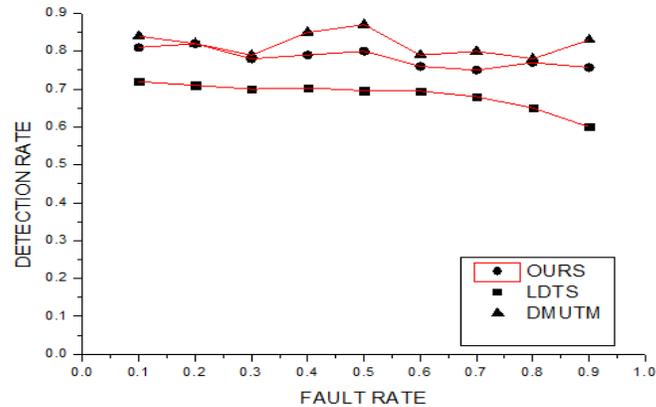


Figure 6. Comparison of Fault Detection Rate

## 6. Conclusions

In this paper, we investigate a method of light-weighted trust calculating. A CM's trust value can be calculated by direct and indirect observation. Direct trust is evaluated by the number of successful and unsuccessful interactions, similar or dissimilar data comparison. Indirect trust is evaluated by aid of similarity matrix in CH. This model can successfully detect fault nodes. We did a series of simulations to test the performance of our proposed model. The simulation results have showed that our model has implemented a balance between detection rate and communication consumption.

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